

Multi-Agent Systems

Kai Cai

cai@omu.ac.jp

Last week

a system of n interacting agents
is modeled by graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$
each agent v_i has an initial value $x_i(0)$

Distributed algorithm

$$x_i(k+1) = x_i(k) + a_{ij} \sum_{j \in \mathcal{N}_i} (x_j(k) - x_i(k))$$
$$(a_{ij} > 0, \sum_{j \in \mathcal{N}_i} a_{ij} < 1)$$

$$x(k+1) = (I - L)x(k) \quad x := \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}_2$$

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a system of n interacting agents

is modeled by graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$

each agent v_i has an initial value $x_i(0)$

$x(k+1) = (I - L)x(k)$ solves averaging

$$\text{i.e. } (\forall v_i \in \mathcal{V}) \lim_{k \rightarrow \infty} x_i(k) = \frac{1}{n} \sum_{v_i \in \mathcal{V}} x_i(0)$$

iff \mathcal{G} is strongly connected and
weight balanced

Last week

a system of n interacting agents
is modeled by graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$
each agent v_i has a local cost function

$$\text{quadratic } f_i(x) = x^2 + \alpha_i x + \beta_i$$

system's global cost function

$$F(x) := \frac{1}{n} \sum_{v_i \in \mathcal{V}} f_i(x)$$

which has a unique optimal solution x^*

Problem: update $x_i(k)$, $k = 1, 2, \dots$,
s.t. $(\forall v_i \in \mathcal{V}) \lim_{k \rightarrow \infty} x_i(k) = x^*$

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Distributed algorithm

$$x(k+1) = (I - L)x(k) \quad x := \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

$$\text{with } x(0) = \left[-\frac{\alpha_1}{2} \quad \dots \quad -\frac{\alpha_n}{2} \right]^\top$$

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is modeled by graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$

each agent v_i has a local cost function

$$\text{quadratic } f_i(x) = x^2 + \alpha_i x + \beta_i$$

$x(k+1) = (I - L)x(k)$ with initial

condition $x(0) = [-\frac{\alpha_1}{2} \ \dots \ -\frac{\alpha_n}{2}]^\top$ solves

optimization , i.e. $(\forall v_i \in \mathcal{V}) \lim_{k \rightarrow \infty} x_i(k) = x^*$

iff \mathcal{G} is strongly connected and
weight balanced

Consensus/ Rendezvous

Examples



Flocking



Schooling



Synchronization



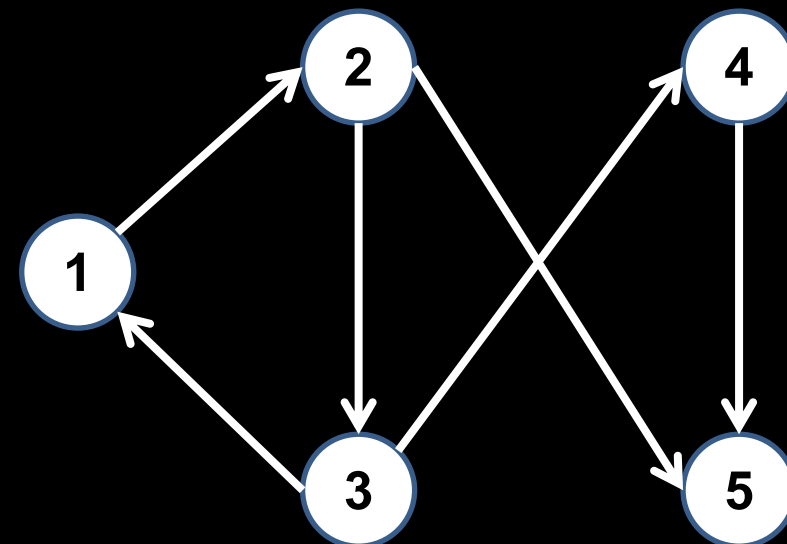
Multi-agent system

a system of n interacting agents
is modeled by graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$

node $v_i \in \mathcal{V}$: an agent

edge $(v_j, v_i) \in \mathcal{E}$: agent j sends
information to v_i

example:



Consensus problem

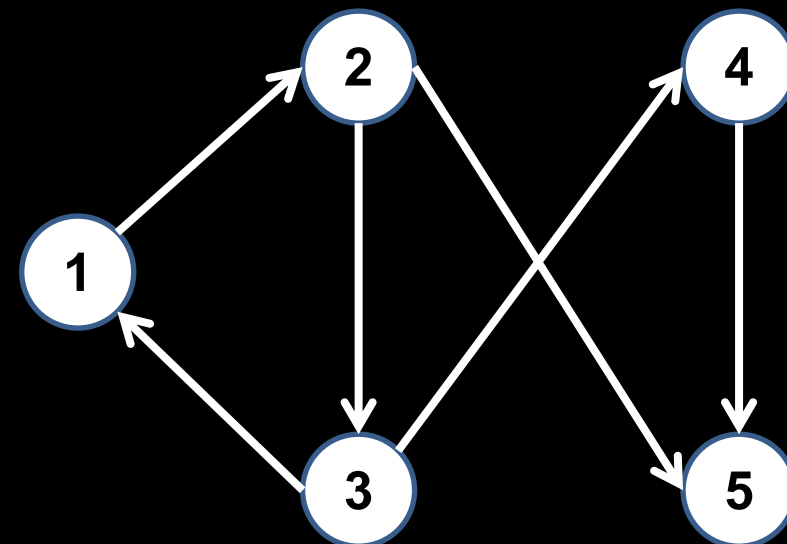
each agent v_i updates its state based on

$$\dot{x}_i = u_i, \quad x_i, u_i \in \mathbb{R}$$

consensus: design input $u_i(t)$, $t \geq 0$

s.t. $(\forall x_1(0), \dots, x_n(0)) (\exists c \in \mathbb{R}) x_i(t) \rightarrow c$
as $t \rightarrow \infty$

example:

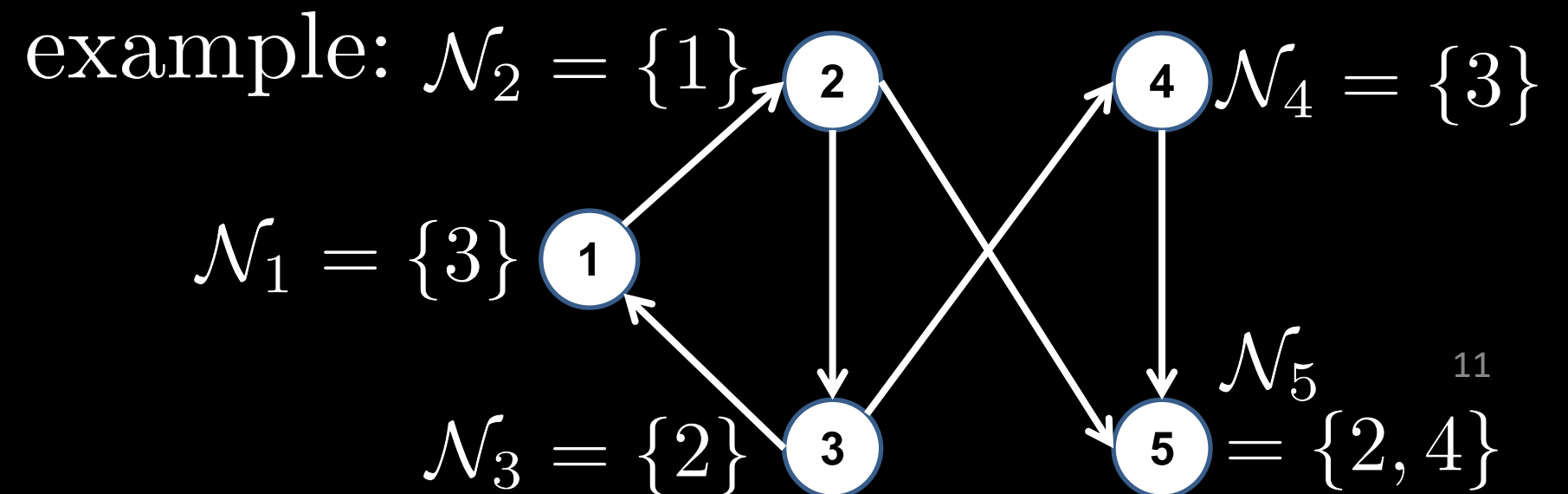


Distributed algorithm

each agent v_i can receive information $x_j(t)$ from neighbor(s) $j \in \mathcal{N}_i$

distributed algorithm: at time $t(\geq 0)$

design $u_i(t)$ based on information $x_j(t)$
where $j \in \mathcal{N}_i$



Example

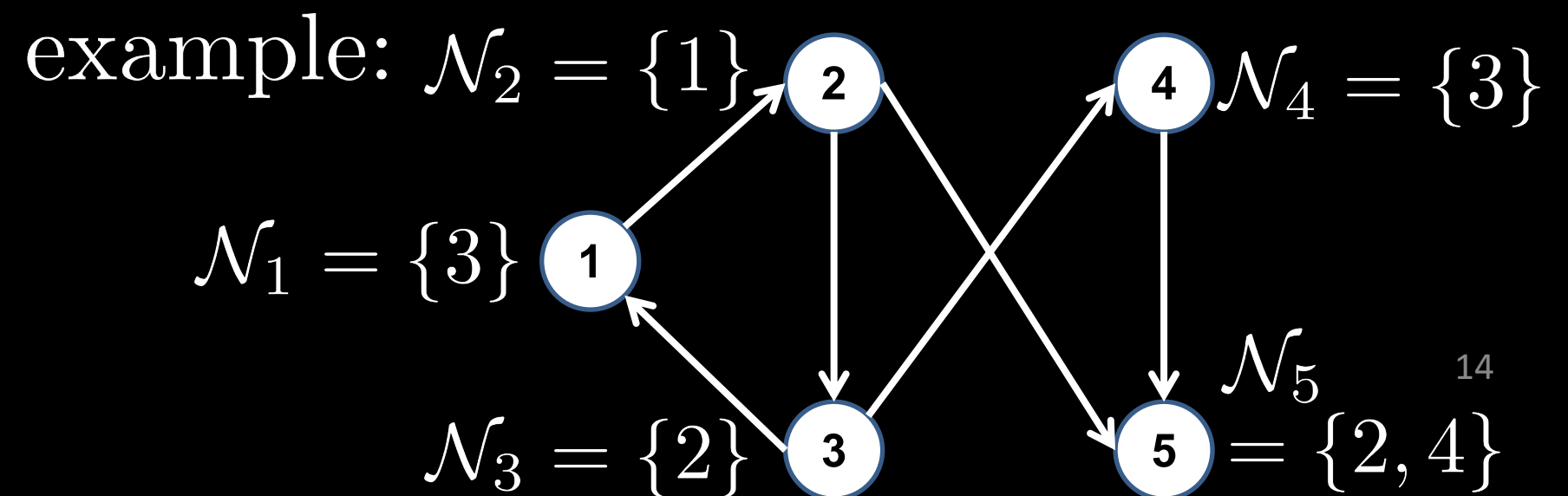
$$\dot{x}_1 = u_1 = x_3 - x_1$$

$$\dot{x}_2 = u_2 = x_1 - x_2$$

$$\dot{x}_3 = u_3 = x_2 - x_3$$

$$\dot{x}_4 = u_4 = x_3 - x_4$$

$$\dot{x}_5 = u_5 = (x_2 - x_5) + (x_4 - x_5)$$



Example

$$\dot{x}_1 = u_1 = x_3 - x_1$$

$$\dot{x}_2 = u_2 = x_1 - x_2$$

$$\dot{x}_3 = u_3 = x_2 - x_3$$

$$\dot{x}_4 = u_4 = x_3 - x_4$$

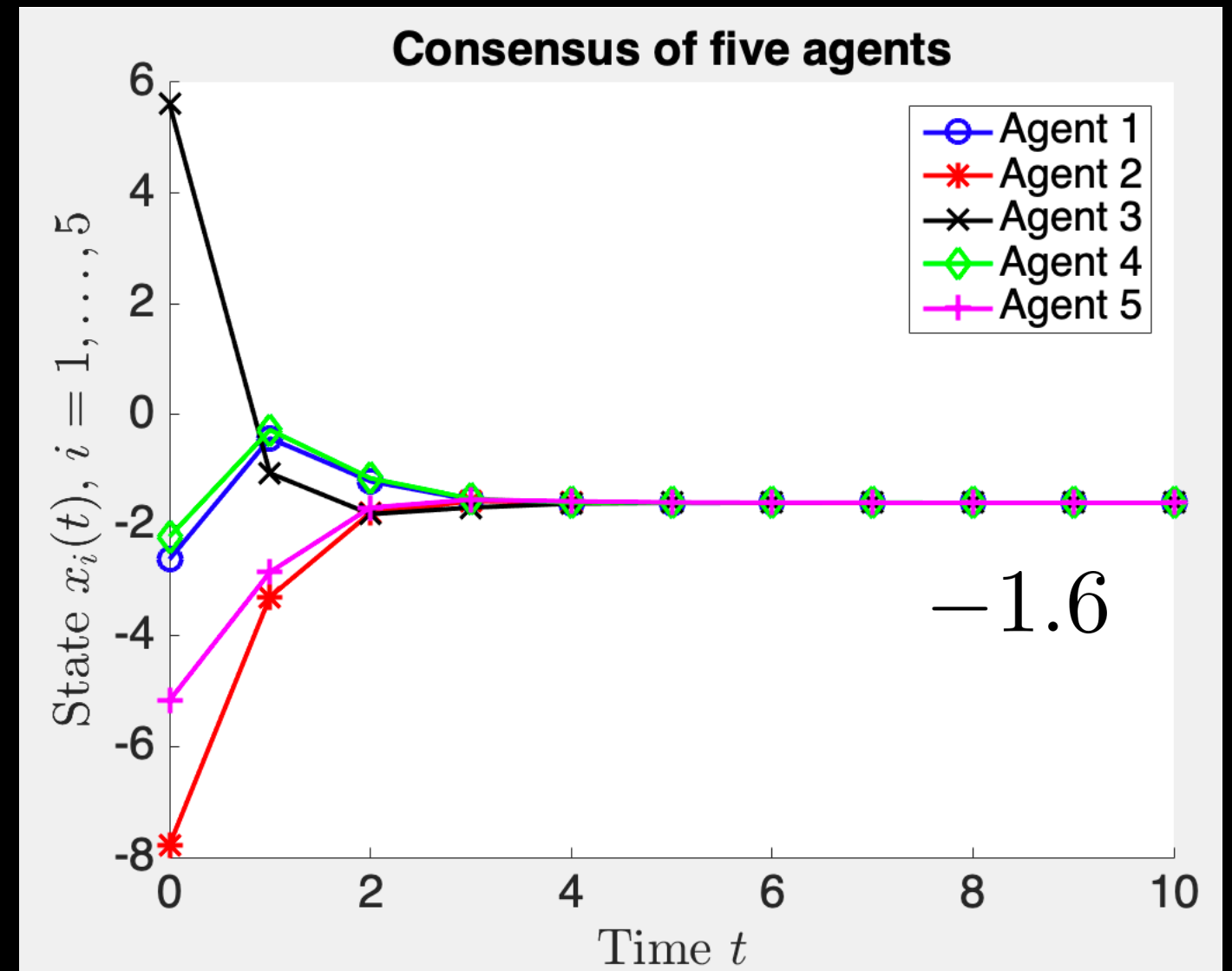
$$\dot{x}_5 = u_5 = (x_2 - x_5) + (x_4 - x_5)$$

$$\dot{x}_i = u_i = \sum_{j \in \mathcal{N}_i} (x_j - x_i)$$

relative state information

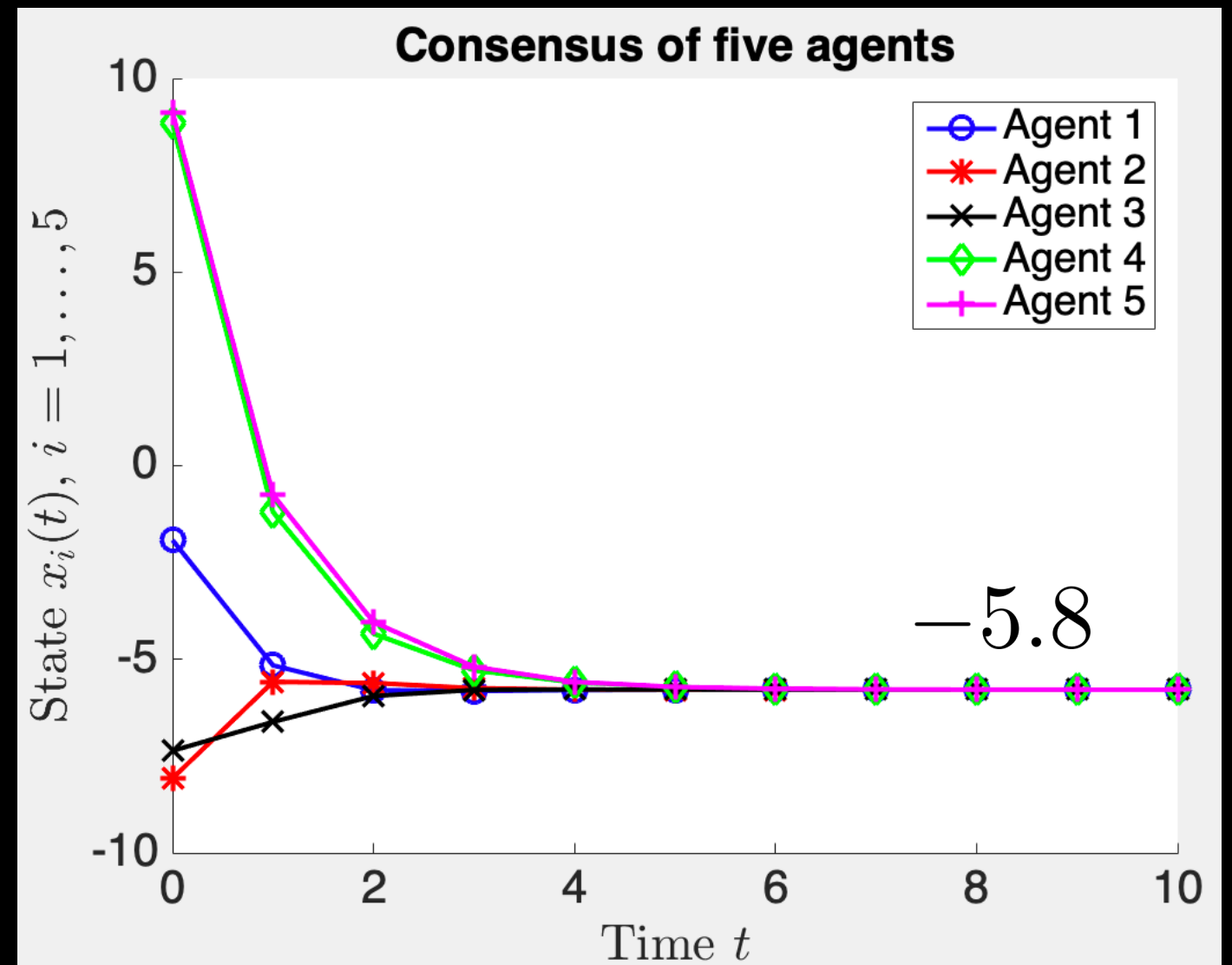
Example

simulation: $x_1(0) = -2.6, x_2(0) = -7.8$
 $x_3(0) = 5.6, x_4(0) = 2.2, x_5(0) = -5.2$



Example

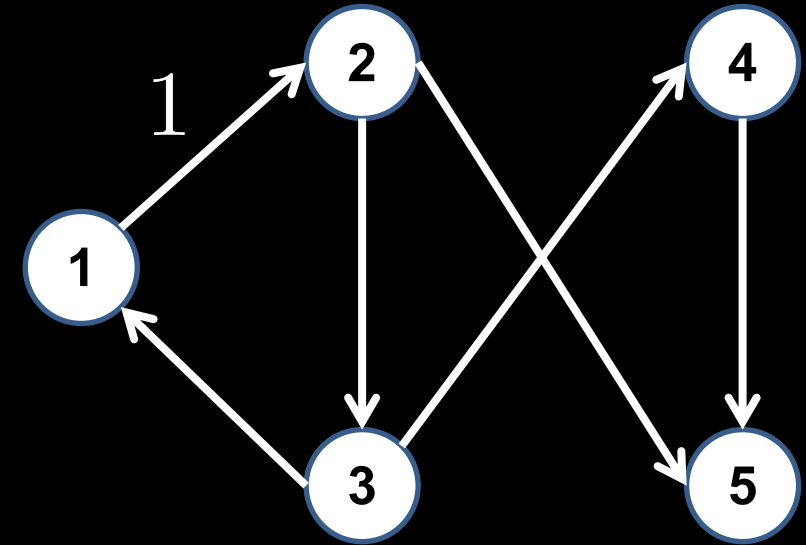
simulation: $x_1(0) = -1.9, x_2(0) = -8.1$
 $x_3(0) = -7.4, x_4(0) = 8.8, x_5(0) = 9.1$



Weighted graph

example:

weighted graph \mathcal{G}



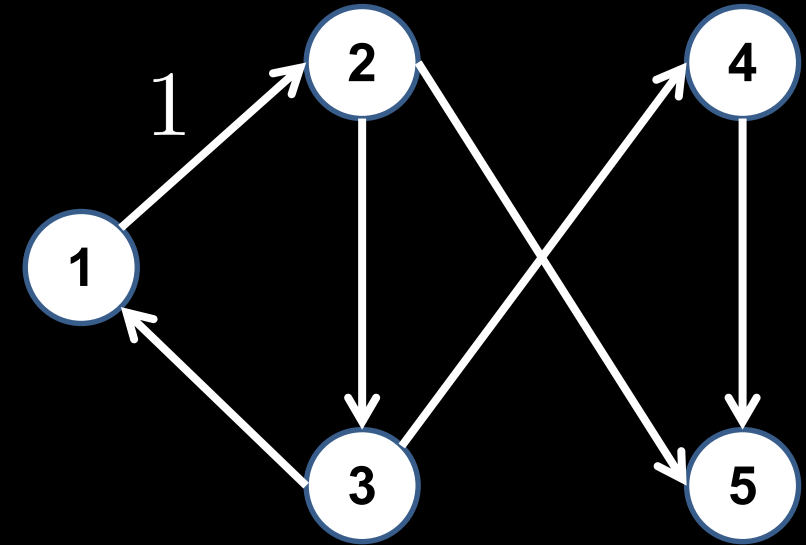
adjacency matrix

$$A = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \end{bmatrix}$$

Weighted graph

example:

weighted graph \mathcal{G}



Laplacian matrix

$$L = \begin{bmatrix} 1 & 0 & -1 & 0 & 0 \\ -1 & 1 & 0 & 0 & 0 \\ 0 & -1 & 1 & 0 & 0 \\ 0 & 0 & -1 & 1 & 0 \\ 0 & -1 & 0 & -1 & 2 \end{bmatrix}$$

Equation

$$\dot{x}_1 = u_1 = x_3 - x_1$$

$$\dot{x}_2 = u_2 = x_1 - x_2$$

$$\dot{x}_3 = u_3 = x_2 - x_3$$

$$\dot{x}_4 = u_4 = x_3 - x_4$$

$$\dot{x}_5 = u_5 = (x_2 - x_5) + (x_4 - x_5)$$

$$\begin{bmatrix} \dot{x}_1 \\ \dot{x}_2 \\ \dot{x}_3 \\ \dot{x}_4 \\ \dot{x}_5 \end{bmatrix} = \begin{bmatrix} -1 & 0 & 1 & 0 & 0 \\ 1 & -1 & 0 & 0 & 0 \\ 0 & 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & -1 & 0 \\ 0 & 1 & 0 & 1 & -2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \end{bmatrix}$$

$-L$

Recap, generalization

a system of n interacting agents

is modeled by graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$

each agent v_i updates its state based on

$$\dot{x}_i = u_i, \quad x_i, u_i \in \mathbb{R}$$

Problem: design u_i to update x_i

$$\text{s.t. } (\forall v_i \in \mathcal{V})(\forall x_i(0))(\exists c) \lim_{t \rightarrow \infty} x_i(t) = c$$

Recap, generalization

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is modeled by graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$

each agent v_i updates its state based on
$$\dot{x}_i = u_i, \quad x_i, u_i \in \mathbb{R}$$

Distributed algorithm

$$\dot{x}_i = u_i = \sum_{j \in \mathcal{N}_i} (x_j - x_i)$$

based on information $x_j(t)$ or
relative information $x_j(t) - x_i(t)$
from neighbor agent(s) $j \in \mathcal{N}_i$

Recap, generalization

a system of n interacting agents
is modeled by graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$

each agent v_i updates its state based on
$$\dot{x}_i = u_i, \quad x_i, u_i \in \mathbb{R}$$

$$x := \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

$$\dot{x} = -Lx$$

Theorem

a system of n interacting agents
is modeled by graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$

each agent v_i updates its state based on
$$\dot{x}_i = u_i, \quad x_i, u_i \in \mathbb{R}$$

$\dot{x} = -Lx$ solves consensus

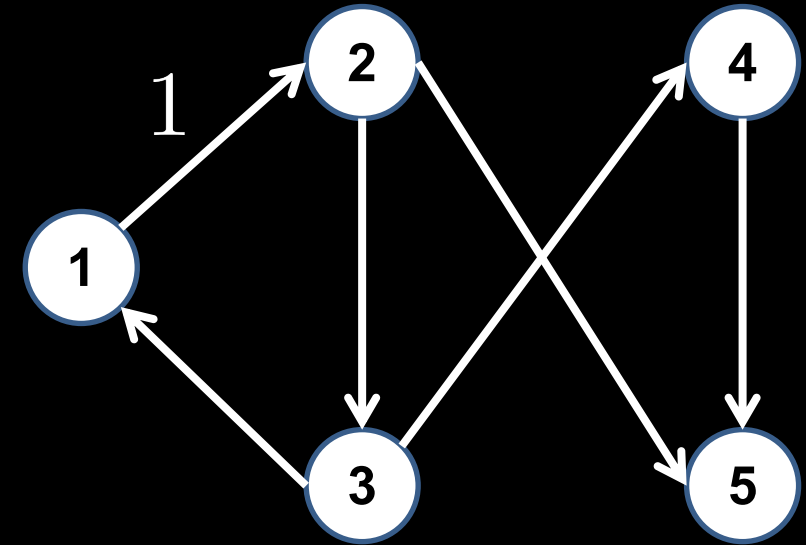
s.t. $(\forall v_i \in \mathcal{V})(\forall x_i(0))(\exists c) \lim_{t \rightarrow \infty} x_i(t) = c$

iff \mathcal{G} contains a spanning tree

Example

example:

weighted graph \mathcal{G}



spanning tree (?)